

ON THE ENUMERATION OF THE SET OF ELEMENTARY  
NUMERICAL SEMIGROUPS WITH FIXED MULTIPLICITY,  
FROBENIUS NUMBER OR GENUS

J. C. ROSALES<sup>1</sup> AND M. B. BRANCO<sup>2</sup>

ABSTRACT. In this paper we give algorithms that allow to compute the set of every elementary numerical semigroups with given genus, Frobenius number and multiplicity. As a consequence we obtain formulas for the cardinality of these sets.

1. INTRODUCTION

Let  $\mathbb{N}$  be the set of nonnegative integers. A numerical semigroup is a subset  $S$  of  $\mathbb{N}$  which is closed under addition,  $0 \in S$  and  $\mathbb{N} \setminus S$  has finitely many elements. The cardinality of the set  $\mathbb{N} \setminus S$  is called the genus of  $S$  and it is denoted by  $g(S)$ .

Given a positive integer  $g$ , we denote by  $\mathcal{S}(g)$  the set of all numerical semigroups with genus  $g$ . The problem of determining the cardinality of  $\mathcal{S}(g)$  has been widely treated in the literature (see for example [2, 4–7] and [13]). Some of these works were motivated by Amorós’s conjecture [5], which says that the sequence of cardinals of  $\mathcal{S}(g)$  for  $g = 1, 2, \dots$  has a Fibonacci behavior. It is still not known in general if for a fixed positive integer  $g$  there are more numerical semigroups with genus  $g + 1$  than numerical semigroups with genus  $g$ .

An algorithm that allows us to compute the set of numerical semigroups with genus  $g$  is provided in [3], where elementary numerical semigroups play an important role. In fact, in [3] an equivalence binary relation  $R$  is defined over  $\mathcal{S}(g)$  such that  $\frac{\mathcal{S}(g)}{R} = \{[S] \mid S \text{ is a elementary numerical semigroup with genus } g\}$ . Moreover, it is proved that if  $S$  and  $T$  are elementary numerical semigroups with genus  $g$  then  $[S] = [T]$  if and only if  $S = T$ . The main idea of the algorithm in [3] is to compute

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every elementary numerical semigroups  $S$  with genus  $g$  and, then, to enumerate the elements in  $[S]$  for each  $S$ .

For any numerical semigroup  $S$ , the smallest positive integer belonging to  $S$  (respectively, the greatest that does not belong to  $S$ ) is called the multiplicity (respectively Frobenius number) of  $S$  and it is denoted by  $m(S)$  (respectively  $F(S)$ ) (see [9]).

We say that a numerical semigroup  $S$  is elementary if  $F(S) < 2m(S)$ . This type of numerical semigroups were also studied in [8] and [13]. We denote by  $\mathcal{E}(m, F, g)$  the set of elementary numerical semigroups with multiplicity  $m$ , Frobenius number  $F$  and genus  $g$  (when one of the parameters to  $\mathcal{E}(m, F, g)$  is replaced by the symbol  $-$ , it represents the set of elementary numerical semigroups in which no restrictions are placed on that parameter).

For any finite set  $A$ ,  $\#A$  denotes the cardinal of  $A$ . Given a rational number  $q$  we denote by  $\lceil q \rceil = \min \{z \in \mathbb{Z} \mid q \leq z\}$ .

In Section 2, we review the results of Y. Zhao in [13] which give formulas for  $\#\mathcal{E}(m, -, -)$ ,  $\#\mathcal{E}(m, -, g)$  and  $\#\mathcal{E}(-, -, g)$ , and state that  $\#\mathcal{E}(-, -, g + 1) = \#\mathcal{E}(-, -, g) + \#\mathcal{E}(-, -, g - 1)$ . Therefore, we get that  $\{\#\mathcal{E}(-, -, g)\}_{g \in \mathbb{N}}$  is a Fibonacci sequence.

In Section 3, we study the sets  $\mathcal{E}(m, F, -)$  and  $\mathcal{E}(-, F, -)$ , find formulas for their cardinality, and describe the behavior of the sequence of cardinals of  $\mathcal{E}(-, F, -)$ .

In Section 4, we present algorithms for calculating  $\mathcal{E}(-, F, g)$  and  $\mathcal{E}(m, F, g)$ . From these algorithms, we can derive the cardinality of these sets.

Finally, in Section 5 we show that the set of all elementary numerical semigroups  $\mathcal{E}$  is a Frobenius variety. This fact, together with the results of [11], allows us to construct recursively the set  $\mathcal{E}$ .

## 2. MULTIPLICITY AND GENUS

Our aim in this section is to see that  $\{\#\mathcal{E}(-, -, g)\}_{g \in \mathbb{N}}$  is a Fibonacci sequence. The next result is easy to prove and appears in [13, Proposition 2.1].

**Lemma 2.1.** *Let  $m$  be an integer such that  $m \geq 2$  and let  $A$  be a subset of  $\{m + 1, \dots, 2m - 1\}$ . Then  $\{0, m\} \cup A \cup \{2m, \rightarrow\}$  is an elementary numerical semigroup with multiplicity  $m$ . Moreover, every elementary numerical semigroup with multiplicity  $m$  is of this form.*

As consequence of the above lemma we have that  $\#\mathcal{E}(m, -, -)$  is equal to the number of subsets of a set with  $m - 1$  elements.

**Corollary 2.1.** *If  $m$  is a positive integer, then  $\#\mathcal{E}(m, -, -) = 2^{m-1}$ .*

The following result is easy to prove and gives conditions imposed on two positive integers  $m$  and  $g$  so that there exists at least one elementary numerical semigroup with multiplicity  $m$  and genus  $g$ .

**Proposition 2.1.** *Let  $m$  and  $g$  be nonnegative integers with  $m \neq 0$ . Then  $\mathcal{E}(m, -, g) \neq \emptyset$  if and only if  $m - 1 \leq g \leq 2(m - 1)$ .*

From Lemma 2.1, we know that  $S \in \mathcal{E}(m, -, g)$  if and only if  $S = \{0, m\} \cup A \cup \{2m, \rightarrow\}$ , where  $A$  is a subset of  $\{m + 1, \dots, 2m - 1\}$  and  $\#A = 2(m - 1) - g$ . So we have the following result, which is also in [13, Corollary 2.2].

**Corollary 2.2.** *Let  $m$  and  $g$  be positive integers such that  $m - 1 \leq g \leq 2(m - 1)$ . Then  $\#\mathcal{E}(m, -, g) = \binom{m-1}{g-(m-1)}$ .*

From the results above we get

$$\mathcal{E}(-, -, g) = \bigcup_{m=\lceil \frac{g}{2} \rceil + 1}^{g+1} \mathcal{E}(m, -, g).$$

Thus we have the following algorithm.

*Algorithm 2.1.* Input:  $g$  positive integer. Output:  $\mathcal{E}(-, -, g)$ .

- 1) For all  $m \in \{\lceil \frac{g}{2} \rceil + 1, \dots, g + 1\}$  compute the set  $\mathcal{E}(m, -, g)$ .
- 2) Return  $\bigcup_{m=\lceil \frac{g}{2} \rceil + 1}^{g+1} \mathcal{E}(m, -, g)$ .

Clearly, we get

$$\#\mathcal{E}(-, -, g) = \sum_{m=\lceil \frac{g}{2} \rceil + 1}^{g+1} \#\mathcal{E}(m, -, g).$$

By applying Corollary 2.2, we obtain the following result.

**Corollary 2.3.** *If  $g$  is a positive integer, then  $\#\mathcal{E}(-, -, g) = \sum_{i=\lceil \frac{g}{2} \rceil}^g \binom{i}{g-i}$ .*

The Fibonacci sequence is the sequence of positive integers defined by the linear recurrence equation  $a_{n+1} = a_n + a_{n-1}$ , with  $a_0 = a_1 = 1$ .

It is clear that  $\mathcal{E}(-, -, 0) = \{\mathbb{N}\}$  and  $\mathcal{E}(-, -, 1) = \{\{0, 2, \rightarrow\}\}$  and so  $\#\mathcal{E}(-, -, 0) = \#\mathcal{E}(-, -, 1) = 1$ . By using Corollary 2.3, we can obtain [13, Proposition 2.3], which states that  $\{\#\mathcal{E}(-, -, g)\}_{g \in \mathbb{N}}$  is a Fibonacci sequence.

**Theorem 2.1.** *If  $g$  is a positive integer, then  $\#\mathcal{E}(-, -, g + 1) = \#\mathcal{E}(-, -, g) + \#\mathcal{E}(-, -, g - 1)$ .*

### 3. MULTIPLICITY AND FROBENIUS NUMBER

Our first goal in this section is to describe sufficient conditions for two positive integers  $m$  and  $F$  so that there exists at least one elementary numerical semigroups with multiplicity  $m$  and Frobenius number  $F$ .

**Lemma 3.1.** *If  $S$  is an elementary numerical semigroup such that  $S \neq \mathbb{N}$ , then  $\frac{F(S)+1}{2} \leq m(S) \leq F(S) + 1$  and  $m(S) \neq F(S)$ .*

*Proof.* Since  $S \neq \mathbb{N}$ , then  $m(S) \geq 2$  and  $m(S) - 1 \notin S$ . Therefore, we have that  $m(S) - 1 \leq F(S)$ . In addition, as  $S$  is an elementary numerical semigroup then  $F(S) < 2(m(S))$  and thus  $F(S) + 1 \leq 2(m(S))$ . □

From the previous lemma we obtain the following result.

**Proposition 3.1.** *Let  $m$  and  $F$  be positive integers. Then  $\mathcal{E}(m, F, -) \neq \emptyset$  if and only if  $\frac{F+1}{2} \leq m \leq F + 1$  and  $m \neq F$ .*

It is clear that  $\mathcal{E}(F + 1, F, -) = \{\{0, F + 1, \rightarrow\}\}$  and  $\mathcal{E}(F - 1, F, -) = \{\{0, F - 1, F + 1, \rightarrow\}\}$ . Hence, we can assume that  $F = m + i$ , where  $i \in \{2, \dots, m - 1\}$ . By applying Lemma 2.1, we deduce that  $S \in \mathcal{E}(m, F, -)$  if and only if there exists  $A \subseteq \{m + 1, \dots, m + i - 1\}$  such that  $S = \{0, m\} \cup A \cup \{F + 1, \rightarrow\}$ . As a consequence we have the following algorithm.

*Algorithm 3.1.* Input:  $m$  and  $F$  positive integers such that  $\frac{F+1}{2} \leq m \leq F + 1$  and  $m \neq F$ .

Output:  $\mathcal{E}(m, F, -)$ .

- 1) If  $m = F + 1$ , then return  $\{\{0, F + 1, \rightarrow\}\}$ .
- 2) If  $m = F - 1$ , then return  $\{\{0, F - 1, F + 1, \rightarrow\}\}$ .
- 3) Compute the set  $C = \{A \mid A \subseteq \{m + 1, \dots, F - 1\}\}$ .
- 4) Return  $\{\{0, m\} \cup A \cup \{F + 1, \rightarrow\} \mid A \in C\}$ .

Gathering all this information, we obtain the following result which can also be deduced from equation (6) of [1].

**Corollary 3.1.** *Let  $m$  and  $F$  be positive integers such that  $\frac{F+1}{2} \leq m \leq F + 1$  and  $m \neq F$ . Then*

$$\#\mathcal{E}(m, F, -) = \begin{cases} 1, & \text{if } m = F + 1, \\ 2^{F-m-1}, & \text{otherwise.} \end{cases}$$

Next we obtain an algorithm that allows us to compute every elementary numerical semigroup with a given Frobenius number. As a consequence of Proposition 3.1, we have

$$\mathcal{E}(-, F, -) = \bigcup_{m \in \{\lceil \frac{F+1}{2} \rceil, \dots, F+1\} \setminus \{F\}} \mathcal{E}(m, F, -).$$

*Algorithm 3.2.* Input:  $F$  positive integer.

Output:  $\mathcal{E}(-, F, -)$ .

- 1) For all  $m \in \{\lceil \frac{F+1}{2} \rceil, \dots, F + 1\} \setminus \{F\}$  compute (using Algorithm 3.1) the set  $\mathcal{E}(m, F, -)$ .
- 2) Return  $\mathcal{E}(-, F, -) = \bigcup_{m \in \{\lceil \frac{F+1}{2} \rceil, \dots, F+1\} \setminus \{F\}} \mathcal{E}(m, F, -)$ .

Therefore, we have  $\#\mathcal{E}(-, F, -) = \sum_{m \in \{\lceil \frac{F+1}{2} \rceil, \dots, F+1\} \setminus \{F\}} \#\mathcal{E}(m, F, -)$ . From Corollary 3.1 we obtain the following result.

**Corollary 3.2.** *If  $F$  is a positive integer, then  $\#\mathcal{E}(-, F, -) = 2^{F - \lceil \frac{F+1}{2} \rceil}$ .*

We finish this section by describing the behavior of the sequence of cardinalities of  $\mathcal{E}(-, F, -)$  for  $F = 1, 2, \dots$ . Observe that  $\#\mathcal{E}(-, 1, -) = \#\mathcal{E}(-, 2, -) = 1$ .

**Proposition 3.2.** *Let  $F$  be an integer greater than or equal to 2.*

- 1) *If  $F$  is odd, then  $\#\mathcal{E}(-, F + 1, -) = \#\mathcal{E}(-, F, -)$ .*
- 2) *If  $F$  is even, then  $\#\mathcal{E}(-, F + 1, -) = \#\mathcal{E}(-, F, -) + \#\mathcal{E}(-, F - 1, -)$ .*

*Proof.* 1) From Corollary 3.2 it is guaranteed that  $\#\mathcal{E}(-, F, -) = 2^{F - \lceil \frac{F+1}{2} \rceil} = 2^{F - \frac{F+1}{2}} = 2^{\frac{F-1}{2}}$ . By repeating this argument we obtain  $\#\mathcal{E}(-, F + 1, -) = 2^{\frac{F-1}{2}}$ . Therefore, we have  $\#\mathcal{E}(-, F + 1, -) = \#\mathcal{E}(-, F, -)$ .

2) Again, by Corollary 3.2, we know that  $\#\mathcal{E}(-, F, -) + \#\mathcal{E}(-, F - 1, -) = 2^{F - \lceil \frac{F+1}{2} \rceil} + 2^{F-1 - \lceil \frac{F}{2} \rceil} = 2^{F - \frac{F+2}{2}} + 2^{F-1 - \frac{F}{2}} = 2^{\frac{F}{2}}$ . We obtain  $\#\mathcal{E}(-, F + 1, -) = 2^{F+1 - \lceil \frac{F+2}{2} \rceil} = 2^{F+1 - \frac{F+2}{2}} = 2^{\frac{F}{2}}$ . Consequently,  $\#\mathcal{E}(-, F + 1, -) = \#\mathcal{E}(-, F, -) + \#\mathcal{E}(-, F - 1, -)$  □

#### 4. MULTIPLICITY, FROBENIUS NUMBER AND GENUS

In this section, we aim to find conditions for  $m$ ,  $F$  and  $g$  positive integers so that there exists at least one elementary numerical semigroup with a given multiplicity  $m$ , Frobenius number  $F$  and genus  $g$ . The next results are a consequence of the results given in [3, Proposition 2 and Corollary 3].

**Lemma 4.1.** *Let  $F$  and  $g$  be two positive integers. Then  $g \leq F \leq 2g - 1$  if and only if  $\mathcal{E}(-, F, g) \neq \emptyset$ .*

**Lemma 4.2.** *Let  $F$  and  $g$  be two positive integers such that  $g \leq F \leq 2g - 1$ , and let  $\mathcal{A}_{F,g} = \{A \mid A \subseteq \{\lceil \frac{F+1}{2} \rceil, \dots, F - 1\}$  and  $\#A = F - g\}$ . Then  $\mathcal{E}(-, F, g) = \{\{0\} \cup A \cup \{F + 1 \rightarrow\} \mid A \in \mathcal{A}_{F,g}\}$ .*

As an immediate consequence of Lemmas 4.1 and 4.2 we have the following algorithm.

*Algorithm 4.1.* Input:  $F$  and  $g$  positive integers such that  $g \leq F \leq 2g - 1$ .

Output:  $\mathcal{E}(-, F, g)$ .

- 1) Compute the set  $C = \{A \mid A \subseteq \{\lceil \frac{F+1}{2} \rceil, \dots, F - 1\}$  and  $\#A = F - g\}$ .
- 2) Return  $\{\{0\} \cup A \cup \{F + 1, \rightarrow\} \mid A \in C\}$ .

As a consequence of the previous algorithm we obtain the following result which also appears in [3, Corollary 4].

**Corollary 4.1.** *If  $F$  and  $g$  are positive integers such that  $g \leq F \leq 2g - 1$ , then  $\#\mathcal{E}(-, F, g) = \binom{\lceil \frac{F}{2} \rceil - 1}{F - g}$ .*

**Lemma 4.3.** *If  $m$ ,  $F$  and  $g$  are three positive integers such that  $m \geq 2$  and  $\mathcal{E}(m, F, g) \neq \emptyset$ , then  $m - 1 \leq g \leq F < 2m$ .*

*Proof.* Since  $\mathcal{E}(m, F, g) \neq \emptyset$ , then  $\mathcal{E}(m, -, g) \neq \emptyset$  and we have that  $m - 1 \leq g$ . From Lemma 4.1, we deduce that  $g \leq F$ . Finally, by Proposition 3.1, we conclude that  $\frac{F+1}{2} \leq m$  and thus  $F < 2m$ . □

Finally, we present the main result of this section.

**Proposition 4.1.** *Let  $m$ ,  $F$  and  $g$  be three positive integers such that  $m \geq 2$ . Then  $\mathcal{E}(m, F, g) \neq \emptyset$  if and only if one of the following conditions holds:*

- 1)  $(m, F, g) = (m, m - 1, m - 1)$ ;
- 2)  $(m, F, g) = (m, F, m)$  and  $m < F < 2m$ ;
- 3)  $m < g < F < 2m$ .

*Proof. Necessity.* If  $\mathcal{E}(m, F, g) \neq \emptyset$  then by applying Lemma 4.3, we deduce that  $m - 1 \leq g \leq F < 2m$ . Assume that  $S \in \mathcal{E}(m, F, g)$ . We distinguish the following four cases.

- a) If  $g = m - 1$ , then  $S = \{0, m, \rightarrow\}$  and so  $F = m - 1$ . Hence,  $(m, F, g) = (m, m - 1, m - 1)$ .
- b) If  $g = m$ , then  $m < F < 2m$  and  $S = \{0, m, \rightarrow\} \setminus \{F\}$ . Whence,  $(m, F, g) = (m, F, m)$  and  $m < F < 2m$ .
- c) If  $g = F$ , then  $S = \{0, F + 1, \rightarrow\}$  and thus  $F + 1 = m$ . Once again we have  $(m, F, g) = (m, m - 1, m - 1)$ .
- d) If  $g \notin \{m - 1, m, F\}$ , then as  $m - 1 \leq g \leq F < 2m$  and we deduce that  $m < g < F < 2m$ .

*Sufficiency.* It is clear that  $\{0, m, \rightarrow\} \in \mathcal{E}(m, m - 1, m - 1)$  and  $\{0, m, \rightarrow\} \setminus \{F\} \in \mathcal{E}(m, F, m)$ . Suppose that  $m < g < F < 2m$ . Let  $A$  be a subset of  $\{m + 1, \dots, F - 1\}$ , with cardinality  $F - g - 1$ . Since  $g(S) = m - 1 + F - 1 - m - 1 + 1 - \#A + 1 = F - 1 - F + g + 1 = g$ , then  $S = \{0, m\} \cup A \cup \{F + 1, \rightarrow\} \in \mathcal{E}(m, F, g)$ .  $\square$

Notice that, by the sufficiency condition of the proof above, we conclude that, if  $m < g < F < 2m$ , knowing an element in  $\mathcal{E}(m, F, g)$  is the same as knowing a subset of  $\{m + 1, \dots, F - 1\}$  with cardinality  $F - g - 1$ . So we have the following algorithm.

*Algorithm 4.2.* Input:  $m, F$  and  $g$  integers such that  $2 \leq m < g < F < 2m$ .

Output:  $\mathcal{E}(m, F, g)$ .

- 1) Compute  $C = \{A \mid A \subseteq \{m + 1, \dots, F - 1\} \text{ and } \#A = F - g - 1\}$ .
- 2) Return  $\{\{0, m\} \cup A \cup \{F + 1, \rightarrow\} \text{ such that } A \in C\}$ .

Clearly  $\#\mathcal{E}(m, m - 1, m - 1) = \#\mathcal{E}(m, F, m) = 1$ . For the remaining cases the following result gives us the cardinality of  $\mathcal{E}(m, F, g)$ .

**Corollary 4.2.** *Let  $m$ ,  $F$  and  $g$  be positive integers such that  $2 \leq m < g < F \leq 2m$ . Then  $\#\mathcal{E}(m, F, g) = \binom{F - m - 1}{F - g - 1}$ .*

*Proof.* As a consequence of Algorithm 4.2 we have that  $S \in \mathcal{E}(m, F, g)$  if and only if there exists  $A \subseteq \{m + 1, \dots, F - 1\}$ , with cardinality  $F - g - 1$  such that  $S = \{0, m\} \cup A \cup \{F + 1, \rightarrow\}$ .  $\square$

We conclude this section by giving an example that illustrates the previous results.

*Example 4.1.* Let us compute  $\mathcal{E}(4, 7, 5)$ . By Corollary 4.2 we have  $\#\mathcal{E}(4, 7, 5) = \binom{7-4-1}{7-5-1} = \binom{2}{1} = 2$ . Now by using Algorithm 4.2, with  $m = 4$ ,  $F = 7$  and  $g = 5$  we can conclude that  $C = \{\{5\}, \{6\}\}$  and  $\mathcal{E}(4, 7, 5) = \{\{0, 4\} \cup \{5\} \cup \{8, \rightarrow\}, \{0, 4\} \cup \{6\} \cup \{8, \rightarrow\}\}$ .

### 5. FROBENIUS VARIETY

A Frobenius variety (see for example [11]) is a nonempty set  $V$  of numerical semigroups fulfilling the following conditions:

- 1) if  $S$  and  $T$  are in  $V$ , then  $S \cap T \in V$ ;
- 2) if  $S$  is in  $V$  and  $S \neq \mathbb{N}$ , then  $S \cup \{F(S)\} \in V$ .

**Proposition 5.1.**  $\mathcal{E} = \{S \mid S \text{ is an elementary numerical semigroup}\}$  is a Frobenius variety.

*Proof.* If  $S$  and  $T$  belong to  $\mathcal{E}$  it is clear that  $S \cap T$  is a numerical semigroup,

$$F(S \cap T) = \max \{F(S), F(T)\}$$

and

$$m(S \cap T) \geq \max \{m(S), m(T)\}.$$

Therefore,  $F(S \cap T) < 2m(S \cap T)$  and thus  $S \cap T \in \mathcal{E}$ .

If  $S$  is an element in  $\mathcal{E}$  and  $S \neq \mathbb{N}$ , then clearly  $\bar{S} = S \cup \{F(S)\}$  is a numerical semigroup such that  $F(\bar{S}) < F(S)$  and  $m(\bar{S})$  is equal to  $m(S)$  or  $F(S)$ . Therefore,  $F(\bar{S}) < 2m(\bar{S})$  and thus  $\bar{S} \in \mathcal{E}$ . □

We define a directed graph  $G(\mathcal{E})$ , with edges pointing from  $T$  to  $S$ , in the following way: the set of vertices is  $\mathcal{E}$  and  $(T, S) \in \mathcal{E} \times \mathcal{E}$  is an edge of  $G(\mathcal{E})$  if and only if  $S \cup \{F(S)\} = T$ .

The goal of this section is to see that  $G(\mathcal{E})$  is a tree with root equal to  $\mathbb{N}$  and to characterize the sons of a vertex. This fact allows us to recursively construct  $G(\mathcal{E})$  and consequently  $\mathcal{E}$ . To this end we need to introduce some concepts and results.

Given a nonempty subset  $A$  of  $\mathbb{N}$  we will denote by  $\langle A \rangle$  the submonoid of  $(\mathbb{N}, +)$  generated by  $A$ , that is,

$$\langle A \rangle = \{\lambda_1 a_1 + \dots + \lambda_n a_n \mid n \in \mathbb{N} \setminus \{0\}, a_i \in A, \lambda_i \in \mathbb{N} \text{ for all } i \in \{1, \dots, n\}\}.$$

It is well known (see for instance [12]) that every numerical semigroup  $S$  is finitely generated, and therefore there exists a finite subset  $A$  of  $\mathbb{N}$  such that  $S = \langle A \rangle$ . Furthermore, we say that  $A$  is a minimal set of generators of  $S$  if no proper subset of  $A$  generates  $S$ . Every numerical semigroup admits an unique minimal set of generators of  $S$  and we denote this set by  $\text{msg}(S)$ . It is well known (see for instance [12]) that  $\text{msg}(S) = (S \setminus \{0\}) \setminus (S \setminus \{0\} + S \setminus \{0\})$  and if  $x \in S$  then  $S \setminus \{x\}$  is a numerical semigroup if and only if  $x \in \text{msg}(S)$ .

As a consequence of [11, Proposition 24 and Theorem 27] we have the following result.

**Theorem 5.1.** *The graph  $G(\mathcal{E})$  is a tree with root  $\mathbb{N}$ . Furthermore, the sons of a vertex  $S$  of  $G(\mathcal{E})$  are in  $\{S \setminus \{x\} \mid x \in \text{msg}(S), x > F(S) \text{ and } S \setminus \{x\} \in \mathcal{E}\}$ .*

The following result is useful to compute the sons of a vertex of  $G(\mathcal{E})$ .

**Proposition 5.2.** *Let  $S$  be an elementary numerical semigroup and  $x \in \text{msg}(S)$  such that  $x > F(S)$ . Then  $S \setminus \{x\}$  is an elementary numerical semigroup if and only if  $x < 2m(S)$ .*

*Proof.* Suppose that  $S = \{0, m(S), \rightarrow\}$ . Then

$$\text{msg}(S) = \{m(S), m(S) + 1, \dots, 2m(S) - 1\}$$

and clearly the result is true. If  $S \neq \{0, m(S), \rightarrow\}$  then  $m(S \setminus \{x\}) = m(S)$  and  $F(S \setminus \{x\}) = x$ . Therefore,  $S \setminus \{x\}$  is elementary numerical semigroup if and only if  $x < 2m(S)$ .  $\square$

We illustrate the above results with the following example.

*Example 5.1.* Let us compute the sons of vertex  $S = \{0, 5, 6, 9, \rightarrow\}$  of  $G(\mathcal{E})$ . We have  $\text{msg}(S) = \{5, 6, 9, 13\}$ ,  $F(S) = 8$  and  $m(S) = 5$ . Whence  $\{x \in \text{msg}(S) \mid F(S) < x < 2m(S)\} = \{9\}$ . Using Theorem 5.1 and Proposition 5.2 we conclude that  $S$  has an unique son  $S \setminus \{9\} = \langle 5, 6, 13, 14 \rangle$ .

Now, we can recursively construct the tree  $G(\mathcal{E})$  starting with  $\mathbb{N}$  and connecting each vertex with their sons. First we construct  $\text{msg}(S \setminus \{x\})$  from  $\text{msg}(S)$ , when  $x$  is a minimal generator of  $S$  greater than  $F(S)$ . It is clear that if  $\text{msg}(S) = \{m, m + 1, \dots, 2m - 1\}$  which is  $S = \{0, m, \rightarrow\}$  then  $\text{msg}(S \setminus \{m\}) = \{m + 1, m + 2, \dots, 2m + 1\}$ . For the remaining cases, we use the following result that appears in [10, Corollary 18].

**Proposition 5.3.** *Let  $S$  be a numerical semigroup with  $\text{msg}(S) = \{n_1, \dots, n_p\}$ . If  $m(S) = n_1 < n_p$  and  $n_p > F(S)$  then*

$$\text{msg}(S \setminus \{n_p\}) = \begin{cases} \{n_1, \dots, n_{p-1}\}, & \text{if exists } i \in \{2, \dots, p-1\} \text{ such that} \\ & n_p + n_1 - n_i \in S, \\ \{n_1, \dots, n_{p-1}, n_p + n_1\}, & \text{otherwise.} \end{cases}$$

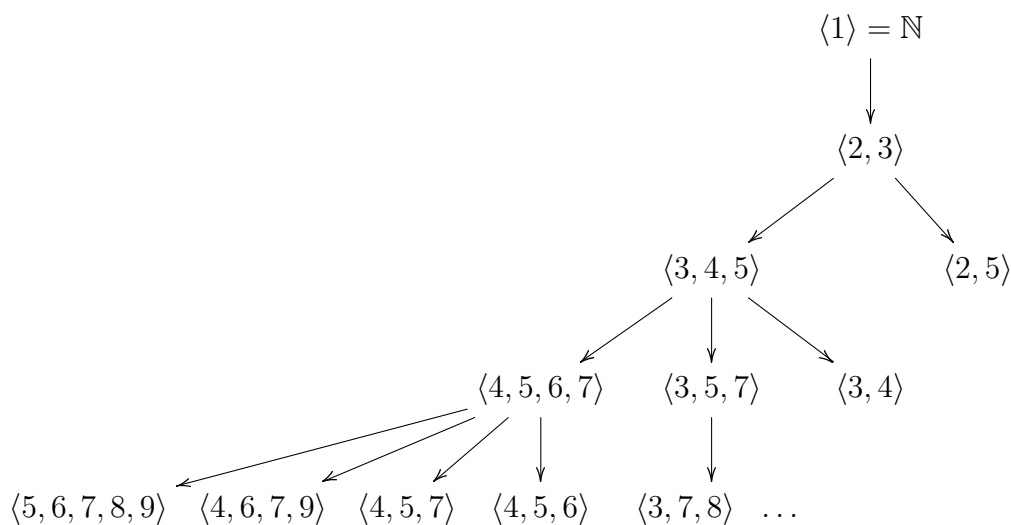
Note that, in the previous proposition, the elements in  $\text{msg}(S)$  are not necessarily ordered.

*Example 5.2.* Let  $S = \langle 5, 6, 9, 13 \rangle$ . Let us compute  $\text{msg}(S \setminus \{9\})$ . By Proposition 5.3, as  $9 + 5 - 6 \notin S$  and  $9 + 5 - 13 \notin S$ , we can conclude that  $\{5, 6, 13, 14\}$  is the minimal system of generators of  $S \setminus \{9\}$ .

Using Theorem 5.1 and Proposition 5.2 and 5.3 we obtain the following:

- .  $\langle 1 \rangle$  has only son  $\langle 1 \rangle \setminus \{1\} = \langle 2, 3 \rangle$ ;
- .  $\langle 2, 3 \rangle$  has two sons  $\langle 2, 3 \rangle \setminus \{2\} = \langle 3, 4, 5 \rangle$  and  $\langle 2, 3 \rangle \setminus \{3\} = \langle 2, 5 \rangle$ ;
- .  $\langle 2, 5 \rangle$  has no sons;





- .  $\langle 3, 4, 5 \rangle$  has three sons  $\langle 3, 4, 5 \rangle \setminus \{3\} = \langle 4, 5, 6, 7 \rangle$ ,  $\langle 3, 4, 5 \rangle \setminus \{4\} = \langle 3, 5, 7 \rangle$  and  $\langle 3, 4, 5 \rangle \setminus \{5\} = \langle 3, 4 \rangle$ ;
- .  $\langle 3, 4 \rangle$  has no sons;
- .  $\langle 3, 5, 7 \rangle$  has one son  $\langle 3, 5, 7 \rangle \setminus \{5\} = \langle 3, 7, 8 \rangle$ ;
- .  $\langle 4, 5, 6, 7 \rangle$  has four sons  $\langle 4, 5, 6, 7 \rangle \setminus \{4\} = \langle 5, 6, 7, 8, 9 \rangle$ ,  $\langle 4, 5, 6, 7 \rangle \setminus \{5\} = \langle 4, 6, 7, 9 \rangle$ ,  $\langle 4, 5, 6, 7 \rangle \setminus \{6\} = \langle 4, 5, 7 \rangle$  and  $\langle 4, 5, 6, 7 \rangle \setminus \{7\} = \langle 4, 5, 6 \rangle$ ;
- . . . . .

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<sup>1</sup>DEPARTAMENTO DE ÁLGEBRA,  
UNIVERSIDAD DE GRANADA,  
E-18071 GRANADA, SPAIN  
*Email address:* `jrosales@ugr.es`

<sup>2</sup>DEPARTAMENTO DE MATEMÁTICA,  
UNIVERSIDADE DE ÉVORA,  
7000 ÉVORA, PORTUGAL  
*Email address:* `mbb@uevora.pt`