




Modularly equidistant numerical semigroups

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Abstract: If S is a numerical semigroup and $s \in S$, we denote by $\text{next}_S(s) = \min \{x \in S \mid s < x\}$. Let a be an integer greater than or equal to two. A numerical semigroup is equidistant modulo a if $\text{next}_S(s) - s - 1$ is a multiple of a for every $s \in S$. In this note, we give algorithms for computing the whole set of equidistant numerical semigroups modulo a with fixed multiplicity, genus, and Frobenius number. Moreover, we will study this kind of semigroups with maximal embedding dimension.

Key words: Embedding dimension, Frobenius number, genus, multiplicity, modularly equidistant numerical semigroups, MED semigroups, numerical semigroup

1. Introduction

Let \mathbb{Z} be the set of integers and let $\mathbb{N} = \{n \in \mathbb{Z} \mid n \geq 0\}$ the set of nonnegative integers. A submonoid of $(\mathbb{N}, +)$ is a subset of \mathbb{N} closed addition and containing 0. A submonoid of $(\mathbb{N}, +)$ with finite complement in \mathbb{N} is a numerical semigroup.

Given a numerical semigroup S and $s \in S$, we denote by $\text{next}_S(s) = \min \{x \in S \mid s < x\}$. For $a \in \mathbb{N} \setminus \{0, 1\}$, we say that S is an equidistant numerical semigroup modulo a if $\text{next}_S(s) - s - 1$ is a multiple of a for every $s \in S$.

We denote by

$$E(a) = \{S \mid S \text{ is equidistant numerical semigroup modulo } a\}.$$

Our aim in this note is to study of this kind of numerical semigroups.

This work is a generalization of the study of the parity numerical semigroups [3]. Indeed a numerical semigroup S is parity if $s + \text{next}_S(s)$ is odd for every $s \in S$. Clearly $s + \text{next}_S(s)$ is odd if and only if $\text{next}_S(s) - s - 1$ is a multiple of 2. Therefore, the parity numerical semigroups are equidistant numerical semigroups modulo 2.

A numerical semigroup S is perfect if $\{x - 1, x + 1\} \subseteq S$ then $x \in S$ (see for instance [4, 5]). Observe that every equidistant modularly numerical semigroup is a perfect numerical semigroup. In fact, if S is equidistant numerical semigroup modulo a and $\{x - 1, x + 1\} \subseteq S$ then $x \in S$, because otherwise $\text{next}_S(x - 1) = x + 1$ and then $\text{next}_S(x - 1) - (x - 1) - 1 = x + 1 - x + 1 - 1 = 1$ that is not a multiple of a .

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We briefly outline the structure of this note. In Section 2, we will order the elements of $E(a)$ in a tree rooted in \mathbb{N} . We will characterize the sons of a vertex and this will allow us to build recursively the elements of $E(a)$.

If S is a numerical semigroup, then $m(S) = \min(S \setminus \{0\})$, $F(S) = \max\{z \in \mathbb{Z} \mid z \notin S\}$ and $g(S) = \text{card}(\mathbb{N} \setminus S)$ (cardinality of $\mathbb{N} \setminus S$) are three important invariants of S called multiplicity, Frobenius number, and genus of S , respectively. See for instance [6] and [2] to understand the importance of the study of these invariants.

As a consequence of the results of Section 2, we present an algorithmic procedure to compute all elements in $E(a)$ with a given multiplicity and genus in Section 3. In Section 4, we get an algorithm for computing the set $E(a)$ with given Frobenius number and we also characterize the maximal elements of this set.

An interesting class of numerical semigroups theory are the numerical semigroups with maximal embedding dimension (see [2]). In Section 5, we characterize the elements in $E(a)$ with maximal embedding dimension.

2. The tree associated to $E(a)$

In this section, we characterize the set $E(a)$, for given a integer greater than or equal to 2. We see how this set can be arranged in a tree.

Lemma 2.1 *If $S \in E(a)$ then $m(S) \equiv 1 \pmod{a}$.*

Proof Clearly, $\text{next}_S(0) = m(S)$ and thus $m(S) - 0 - 1$ is a multiple of a . Therefore, we have that $m(S) \equiv 1 \pmod{a}$. □

It is clear that \mathbb{N} is a numerical semigroup that belongs to $E(a)$. If $S \in E(a)$ and $S \neq \mathbb{N}$ then $m(S) \geq 2$, and so there exists $k \in \mathbb{N} \setminus \{0\}$ such that $m(S) = 1 + ka$. Therefore, $a < m(S)$.

Lemma 2.2 *If $S \in E(a) \setminus \{\mathbb{N}\}$, then $\{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} \subseteq \mathbb{N} \setminus S$.*

Proof Since $F(S) \geq m(S) - 1 \geq a$ we have that $\{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} \subseteq \mathbb{N} \setminus \{0\}$. Assume that there exists $i \in \{1, \dots, a - 1\}$ such that $F(S) - i \in S$ and let $t = \min\{i \in \{1, \dots, a - 1\} \mid F(S) - i \in S\}$. Then $\text{next}_S(F(S) - t) = F(S) + 1$. As $S \in E(a)$ then $F(S) + 1 - (F(S) - t) - 1$ is a multiple of a . Hence, $t = la$ for some $l \in \mathbb{N} \setminus \{0\}$ and consequently $t \geq a$, which is impossible. □

Lemma 2.3 *If $S \in E(a) \setminus \{\mathbb{N}\}$, then $S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} \in E(a)$.*

Proof Since $a < m(S)$, by applying Lemma 2.2, we get that $T = S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\}$ is a numerical semigroup. Let $s \in S$ such that $\text{next}_S(s) = F(S) + 1$. In order to conclude the proof, it suffices to show that $\text{next}_T(s) - s - 1$ is a multiple of a . As $\text{next}_T(s) = F(S) - (a - 1)$ then $\text{next}_T(s) - s - 1$ is multiple of a if and only if $F(S) - (a - 1) - s - 1 = F(S) - s - a$ is multiple of a . However, this is true, because if $S \in E(a)$ then $F(S) + 1 - s - 1 = \text{next}_S(s) - s - 1$ is multiple of a . □

The previous result can be viewed as a procedure to construct a sequence $\{S_n \mid n \in \mathbb{N}\}$ of elements in $E(a)$. Given a numerical semigroup S , we define recursively

- $S_0 = S$,

$$\bullet S_{n+1} = \begin{cases} S_n \cup \{F(S_n), F(S_n) - 1, \dots, F(S_n) - (a - 1)\} & \text{if } S_n \neq \mathbb{N} \\ \mathbb{N} & \text{otherwise.} \end{cases}$$

The next result can be easily proved.

Proposition 2.4 *If $S \in E(a)$ and $\{S_n \mid n \in \mathbb{N}\}$ is the previous sequence, then there exists $k \in \mathbb{N}$ such that $S_k = \mathbb{N}$. Furthermore, for every $i \in \{0, 1, \dots, k - 1\}$ then $\text{card}(S_{i+1} \setminus S_i) = a$ and $k = \frac{g(S)}{a}$.*

We illustrate the previous result with an example.

Example 2.5 *In the following sequence of elements of $E(2)$, we have $k = 3$.*

$$S_0 = \langle 5, 6, 7 \rangle \subsetneq S_1 = \langle 5, 6, 7, 8, 9 \rangle \subsetneq S_2 = \langle 3, 4, 5 \rangle \subsetneq S_3 = \mathbb{N}.$$

A graph G is a pair (V, E) , where V is a nonempty set and E is a subset of $\{(v, w) \in V \times V \mid v \neq w\}$. The elements of V and E are called vertices and edges of G , respectively. A path of length n connecting the vertices u and v of G is a sequence of distinct edges of the form $(v_0, v_1), (v_1, v_2), \dots, (v_{n-1}, v_n)$ with $v_0 = u$ and $v_n = v$.

A graph G is a tree if there exists a vertex r (known as the root of G) such that for every other vertex v of G , there exists a path connecting v and r . If (u, v) is a edge of the tree then we say that u is a son of v .

We define the graph $G(E(a))$ as graph whose vertices are the elements of $E(a)$ and $(S, T) \in E(a) \times E(a)$ is an edge if $T = S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\}$. From Proposition 2.4, we deduce the following.

Theorem 2.6 *The graph $G(E(a))$ is a tree with root equal to \mathbb{N} .*

Note that the tree $G(E(a))$ can be constructed recursively, from the root \mathbb{N} in each step we are joining each of the vertices with its sons. Our next goal is to characterize the sons of an arbitrary vertex in the tree $G(E(a))$. We distinguish two cases depending on whether or not the vertex is \mathbb{N} .

Lemma 2.7 *The vertex \mathbb{N} has a unique son in the tree $G(E(a))$ that is $\{0, a + 1, \rightarrow\}$.*

Proof If S is a son of \mathbb{N} , then $S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} = \mathbb{N}$ and thus $S = \{0, a + 1, \rightarrow\}$. \square

If \mathcal{X} is a nonempty subset of \mathbb{N} , then we will write $\langle \mathcal{X} \rangle$ for the submonoid of $(\mathbb{N}, +)$ generated by \mathcal{X} , that is,

$$\langle \mathcal{X} \rangle := \left\{ \sum_{i=1}^n \lambda_i x_i \mid n \in \mathbb{N} \setminus \{0\}, x_1, \dots, x_n \in \mathcal{X}, \text{ and } \lambda_1, \dots, \lambda_n \in \mathbb{N} \right\}.$$

It is well known (see for example [8]) that $\langle \mathcal{X} \rangle$ is a numerical semigroup if and only if $\text{gcd}(\mathcal{X}) = 1$.

If M is a submonoid of $(\mathbb{N}, +)$ and $M = \langle \mathcal{X} \rangle$ then we say that \mathcal{X} is a system of generators of M . Moreover, if $M \neq \langle \mathcal{Y} \rangle$ for all $\mathcal{Y} \subsetneq \mathcal{X}$, then we say that \mathcal{X} is a minimal system of generators of S . The following result can be deduced from [8, Corollary 2.8].

Corollary 2.8 *If M is a submonoid of $(\mathbb{N}, +)$, then M has a unique minimal system of generators, which in addition is finite.*

This minimal system of generators of M will be denoted by $\text{msg}(M)$.

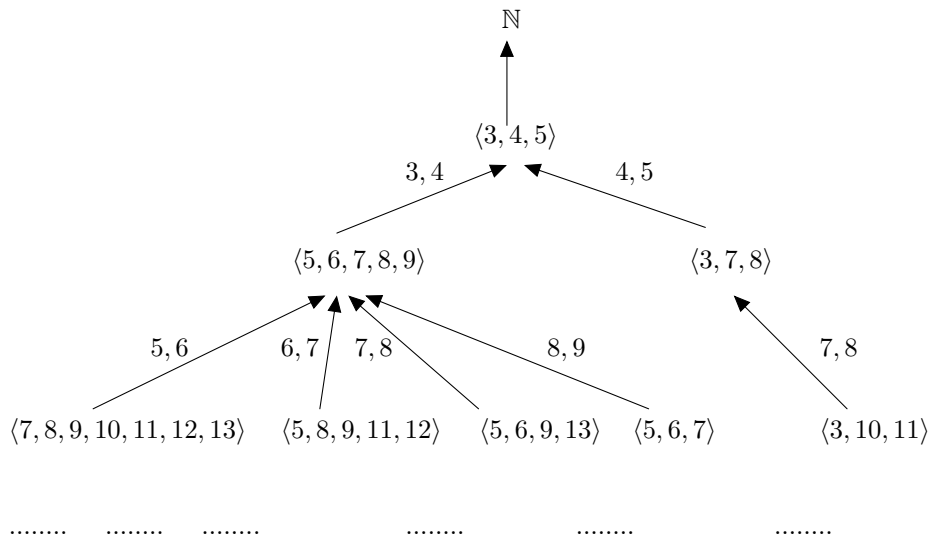
Lemma 2.9 [7, Lemma 1.7] *Let S be a numerical semigroup and $x \in S$. Then $S \setminus \{x\}$ is a numerical semigroup if and only if $x \in \text{msg}(S)$.*

Proposition 2.10 *Let $T \in E(a) \setminus \{\mathbb{N}\}$. Then the set of sons of T , in the tree $G(E(a))$, is equal to $\{T \setminus \{x, x + 1, \dots, x + (a - 1)\} \mid \{x, x + 1, \dots, x + (a - 1)\} \subseteq \text{msg}(T) \text{ and } x > F(T)\}$.*

Proof If $T \in E(a) \setminus \{\mathbb{N}\}$, then by Lemma 2.1, we have that $m(T) > a$. Since S is a son of T then $T = S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\}$. Hence, $\{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} \subseteq \text{msg}(T)$ and $F(S) - (a - 1) > F(T)$.

Conversely, if $\{x, x + 1, \dots, x + (a - 1)\} \subseteq \text{msg}(T)$ and $x > F(T)$ then, using repeatedly Lemma 2.9, we obtain that $S = T \setminus \{x, x + 1, \dots, x + (a - 1)\}$ is a numerical semigroup with $F(S) = x + a - 1$. Therefore, $T = S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\}$ and thus S is a son of T . □

Example 2.11 *Let us construct recursively the tree $G(E(2))$.*



The numbers that appear on either side of the edges is the elements that we remove from the semigroup to obtain its son. For example $\langle 3, 4, 5 \rangle \setminus \{3, 4\} = \langle 5, 6, 7, 8, 9 \rangle$.

Observe that $E(a)$ has infinite cardinality, because $\{0, ka + 1, \rightarrow\} \in E(a)$ for all $k \in \mathbb{N}$.

3. The set $E(a)$ with a given multiplicity and genus

We will denote by $E(a, m) = \{S \in E(a) \mid m(S) = m\}$. By Lemma 2.1, we obtain that $E(a, m) \neq \emptyset$ if and only if $m = ka + 1$ for some $k \in \mathbb{N}$. It is clear that $E(a, 1) = \{\mathbb{N}\}$.

From now on, assume that $m = ka + 1$ with $k \in \mathbb{N} \setminus \{0\}$. It is clear that maximum element in $E(a)$ is $\langle m, m + 1, \dots, 2m - 1 \rangle = \{0, m, \rightarrow\}$. From Lemma 2.3, we deduce the next result.

Lemma 3.1 *If $S \in E(a, m)$ and $S \neq \{0, m, \rightarrow\}$, then $S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} \in E(a, m)$.*

The previous lemma allows us to define recurrently the sequence $\{S_n \mid n \in \mathbb{N}\}$ of elements of $S \in E(a, m)$. If $S \in E(a, m)$, then

- $S_0 = S$,
- $S_{n+1} = \begin{cases} S_n \cup \{F(S_n), F(S_n) - 1, \dots, F(S_n) - (a - 1)\} & \text{if } S_n \neq \{0, m, \rightarrow\} \\ \{0, m, \rightarrow\} & \text{otherwise.} \end{cases}$

The next result has immediate proof.

Proposition 3.2 *If $S \in E(a, m)$ and $\{S_n \mid n \in \mathbb{N}\}$ is the previous sequence, then there exists $k \in \mathbb{N}$ such that $S_k = \{0, m, \rightarrow\}$. Moreover, for every $i \in \{0, 1, \dots, k - 1\}$ then $\text{card}(S_{i+1} \setminus S_i) = a$ and $k = \frac{g(S) - (m - 1)}{a}$.*

As a consequence of Lemma 2.1 and Proposition 3.2, we have the following result.

Corollary 3.3 *If $S \in E(a, m)$ then $g(S)$ is a multiple of a .*

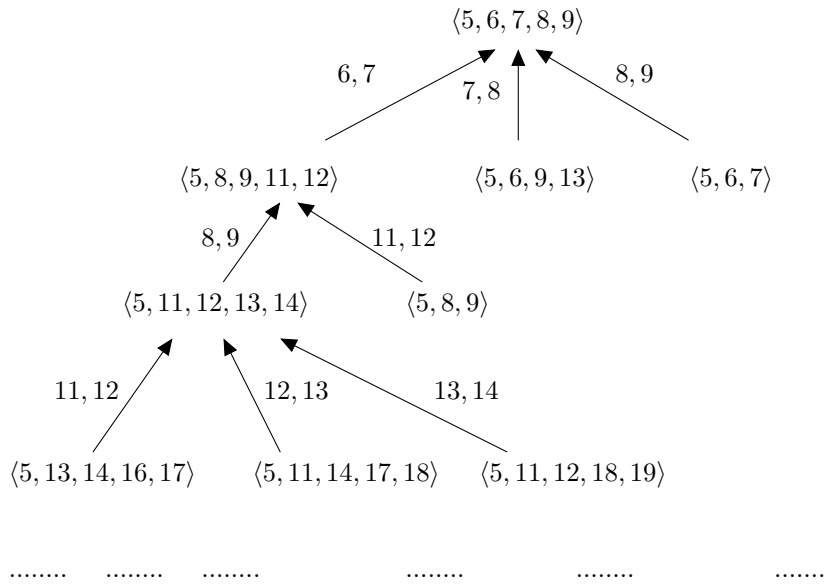
We define the graph $G(E(a, m))$ as follows: $E(a, m)$ is its set of vertices and $(S, T) \in E(a) \times E(a)$ is an edge if $T = S \cup \{F(S), F(S) - 1, \dots, F(S) - (a - 1)\}$.

The following result is a consequence of Propositions 2.10 and 3.2.

Theorem 3.4 *The graph $G(E(a, m))$ is a tree rooted in $\{0, m, \rightarrow\}$. Moreover, the set of sons of T is equal to*

$$\{T \setminus \{x, x + 1, \dots, x + (a - 1)\} \mid \{x, x + 1, \dots, x + (a - 1)\} \subseteq \text{msg}(T), x > F(T) \text{ and } x \neq m\}.$$

Example 3.5 *We are going to build the tree $G(E(2, 5))$.*



Observe that for $t \in \mathbb{N} \setminus \{0\}$ we have that $\langle m \rangle \cup \{tm, \rightarrow\} \in E(a, m)$ and so $E(a, m)$ has infinite cardinality. Our next aim in this section will be to show an algorithm that allows us to compute the set $E(a, m)$ with given genus.

If $G = (V, E)$ is a tree and $v \in V$, then the depth of v , which we will denote by $d(v)$, is the length of a unique path that connects v with the root. Given k a nonnegative integer, denote by $N(G, k) = \{v \in V \mid d(v) = k\}$.

As a consequence of Proposition 3.2 we have the following result.

Proposition 3.6 *Let $S \in E(a, m)$. Then $S \in N(G(E(a, m)), k)$ if and only if $g(S) = m - 1 + ak$.*

The next lemma follows immediately from the definitions.

Lemma 3.7 *If $k \in \mathbb{N}$, then $N(G(E(a, m)), k + 1) = \{S \mid S \text{ is a son of an element in } N(G(E(a, m)), k)\}$.*

We are already able to give an algorithm that allows us to compute the set $E(a, m)$ with a given genus. Note that $S \in E(a, m)$ then $\{1, \dots, m - 1\} \subseteq \mathbb{N} \setminus S$ and thus $g(S) \geq m - 1$. Moreover, by Corollary 3.3, we have that $g(S)$ is a multiple of a .

Algorithm 1

INPUT: a, m and g nonnegative integers such that $2 \leq a \leq m - 1 \leq g$,
 $m \equiv 1 \pmod{a}$ and $g \equiv 0 \pmod{a}$.

OUTPUT: The set $\{S \in E(a, m) \mid g(S) = g\}$.

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1: Set  $i = m - 1$  and  $A = \{m, m + 1, \dots, 2m - 1\}$ 
2: while True do
3:   if  $i = g$  then
4:     return  $A$ 
5:   for  $S \in A$  do
6:      $B_S = \{\{x, x + 1, \dots, x + (a - 1)\} \subseteq \text{msg}(S), x > F(S) \text{ and } x \neq m\}$ 
7:    $A := \bigcup_{S \in A} \{S \setminus \{x, x + 1, \dots, x + (a - 1)\} \mid \{x, x + 1, \dots, x + (a - 1)\} \in B_S\}$ ,
    $i = i + a$  and go to step 2.

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Example 3.8 Let us compute the set $\{S \in E(2, 5) \mid g(S) = 8\}$ using Algorithm 1.

1. Set $i = 4$ and $A = \{5, 6, 7, 8, 9\}$.
2. The first loop constructs $B_{\langle 5,6,7,8,9 \rangle} = \{\{6, 7\}, \{7, 8\}, \{8, 9\}\}$ then
 $A = \{5, 6, 7, 8, 9\} \setminus \{6, 7\}, \{5, 6, 7, 8, 9\} \setminus \{7, 8\}, \{5, 6, 7, 8, 9\} \setminus \{8, 9\}$, $i = 6$,
3. the second loop constructs $B_{\langle 5,8,9,11,12 \rangle} = \{\{8, 9\}, \{11, 12\}\}$, $B_{\langle 5,6,9,13 \rangle} = \emptyset$ and $B_{\langle 5,6,7 \rangle} = \emptyset$ then
 $A = \{5, 8, 9, 11, 12\} \setminus \{8, 9\}, \{5, 8, 9, 11, 12\} \setminus \{11, 12\}$, $i = 8$.

Hence, $\{S \in E(2, 5) \mid g(S) = 8\} = \{5, 11, 12, 13, 14\}, \{5, 8, 9\}$.

4. The set $E(a)$ with a given Frobenius number

Our aim in this section will be to show an algorithm that allows us to compute the set $E(a)$ with a given Frobenius number. If S is a numerical semigroup such that $S \neq \mathbb{N}$, then we have that $F(S) \geq m(S) - 1$. Besides, by Lemma 2.1, if $S \in E(a) \setminus \{\mathbb{N}\}$ then $m(S) = 1 + ka$ for some $k \in \mathbb{N} \setminus \{0\}$. Since $m(S) \leq F(S) + 1$ and $m(S) = 1 + ka$ then $k \leq \frac{F(S)}{a}$.

Given a rational number q we denote by $\lfloor q \rfloor$ its integer part, that is, $\lfloor q \rfloor = \max\{z \in \mathbb{Z} \mid z \leq q\}$. We can announce the following result.

Lemma 4.1 Let $F(S)$ be an integer greater than or equal to two. Then

$$\{S \in E(a) \mid F(S) = F\} = \bigcup_{k=1}^{\lfloor \frac{F}{a} \rfloor} \{S \in E(a, ka + 1) \mid F(S) = F\}.$$

Clearly, $\{S \in E(a, ka + 1) \mid F(S) = ka\} = \{0, ka + 1 \rightarrow\}$.

Lemma 4.2 Let $k \in \mathbb{N} \setminus \{0\}$ such that $ka + 1 < F$. Then $\{S \in E(a, ka + 1) \mid F(S) = F\} \neq \emptyset$ if and only if $F \bmod (ka + 1) \notin \{0, 1, \dots, a - 1\}$.

Proof *Necessity.* If $S \in E(a, ka + 1)$ and $F(S) = F$ then, applying Lemma 2.2, we have that $\{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} \subseteq \mathbb{N} \setminus S$. Therefore, $\{F(S), F(S) - 1, \dots, F(S) - (a - 1)\} \cap \langle ka + 1 \rangle = \emptyset$ and so $F \bmod (ka + 1) \notin \{0, 1, \dots, a - 1\}$.

Sufficiency. Assume that $r = F \bmod (ka + 1)$. Then $F = q(ka + 1) + r$ for some $q \in \mathbb{N} \setminus \{0\}$ and $r \in \{a, \dots, ka\}$. Hence we obtain that

$$S = \{0, ka + 1, 2(ka + 1), \dots, q(ka + 1), q(ka + 1) + 1, \dots, q(ka + 1) + r - a, F + 1, \rightarrow\} \in E(a)$$

and $F(S) = F$. □

Now by using Lemmas 4.1 and 4.2, in order to compute the set $E(a)$ with a given Frobenius number F , it is enough to give an algorithm that computes this set with m a positive integer and verifies that $2 \leq a \leq m - 1$, $m < F$, $m \equiv 1 \pmod{a}$, and $(F \bmod m) \notin \{0, 1, \dots, a - 1\}$.

Algorithm 2

INPUT: a, m and F nonnegative integers such that $2 \leq a \leq m - 1$,
 $m < F$, $m \equiv 1 \pmod{a}$ and $(F \bmod m) \notin \{0, 1, \dots, a - 1\}$.

OUTPUT: The set $\{S \in E(a, m) \mid F(S) = F\}$.

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1:  $B = \emptyset$  and  $A = \{m, m + 1, \dots, 2m - 1\}$ 
2: while True do
3:   for  $S \in A$  do
4:     Compute  $B_S = \{\{x, x + 1, \dots, x + (a - 1)\} \subseteq \text{msg}(S) \mid x \neq m,$ 
 $x > F(S)$  and  $x + (a - 1) \leq F\}$ 
5:    $B := B \cup \{S \setminus \{x, x + 1, \dots, x + (a - 1)\} \mid S \in A,$ 
 $\{x, x + 1, \dots, x + (a - 1)\} \in B_S$  and  $x + (a - 1) = F\}$ 
6:    $A := \bigcup_{S \in A} \{S \setminus \{x, x + 1, \dots, x + (a - 1)\} \mid$ 
 $\{x, x + 1, \dots, x + (a - 1)\} \in B_S$  and  $x + (a - 1) < F\}$ 
7:   if  $A = \emptyset$  then
8:     return  $B$ 

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Next we illustrate this method with an example.

Example 4.3 Let us compute the set $\{S \in E(2) \mid F(S) = 12\}$. First, by using Lemma 4.1, we have that $\{S \in E(2) \mid F(S) = 12\} = \bigcup_{m \in \{3, 5, 7, 9, 11, 13\}} \{S \in E(2, m) \mid F(S) = 12\}$. From Lemma 4.2, we obtain that $\{S \in E(2, 3) \mid F(S) = 12\} = \emptyset$ and $\{S \in E(2, 11) \mid F(S) = 12\} = \emptyset$. Moreover, by the observation made after the Lemma 4.1, we get that $\{S \in E(2, 13) \mid F(S) = 12\} = \{0, 13, \rightarrow\}$. Therefore, by using Algorithm 2, we have to compute the set $\{S \in E(2, m) \mid F(S) = 12\}$ with $m \in \{5, 7, 9\}$.

For example, we will calculate the set $\{S \in E(2, 5) \mid F(S) = 12\}$.

1. Start $B = \emptyset$ and $A = \{5, 6, 7, 8, 9\}$.
2. The first loop constructs $B_{\langle 5, 6, 7, 8, 9 \rangle} = \{\{6, 7\}, \{7, 8\}, \{8, 9\}\}$
then $B = \emptyset$,
3. next constructs $A = \{\langle 5, 6, 7, 8, 9 \rangle \setminus \{6, 7\}, \langle 5, 6, 7, 8, 9 \rangle \setminus \{7, 8\}, \langle 5, 6, 7, 8, 9 \rangle \setminus \{8, 9\}\}$,
4. the second loop constructs $B_{\langle 5, 8, 9, 11, 12 \rangle} = \{\{8, 9\}, \{11, 12\}\}$, $B_{\langle 5, 6, 9, 13 \rangle} = \emptyset$ and $B_{\langle 5, 6, 7 \rangle} = \emptyset$
then $B = \{\langle 5, 8, 9, 11, 12 \rangle \setminus \{11, 12\}\}$,
5. next constructs $A = \{\langle 5, 8, 9, 11, 12 \rangle \setminus \{8, 9\}\}$,

6. the third loop constructs $B_{\langle 5,11,12,13,14 \rangle} = \{\{11, 12\}\}$
then $B = \{\langle 5, 8, 9 \rangle, \langle 5, 11, 12, 13, 14 \rangle \setminus \{11, 12\}\}$,
7. next constructs $A = \emptyset$,
8. $\{S \in E(2, 5) \mid F(S) = 12\} = \{\langle 5, 8, 9 \rangle, \langle 5, 13, 14, 16, 17 \rangle\}$.

Next we are interested in characterizing of the maximal elements in the set $\{S \in E(a) \mid F(S) = F\}$. The next result is well known.

Lemma 4.4 [9, Lemma 10] *Let S and T be two numerical semigroups such that $S \subsetneq T$ and $x = \max(T \setminus S)$. Then $S \cup \{x\}$ is a numerical semigroup.*

Proposition 4.5 *Let $\{S, T\} \subseteq E(a)$ such that $S \subsetneq T$, $x = \max(T \setminus S)$ and let $s \in S$ such that $s < x < \text{next}_S(s)$. If $\{x_1 < x_2 < \dots < x_r\} = \{t \in T \mid s < t < \text{next}_S(s)\}$ then $S \cup \{x_1, x_2, \dots, x_r\} \in E(a)$, r is a multiple of a and $S \cup \{x_r, x_{r-1}, \dots, x_{r-(a-1)}\} \in E(a)$.*

Proof By repeatedly applying Lemma 4.4, we get that $S \cup \{x_1, x_2, \dots, x_r\}$ is a numerical semigroup. Moreover, as $\{S, T\} \subseteq E(a)$ such that $S \subseteq T$, we deduce that $S \cup \{x_1, x_2, \dots, x_r\} \in E(a)$.

Since $s < x_1 < \dots < x_r < \text{next}_S(s)$ are consecutive elements of T and $T \in E(a)$, then there exist $\{k_1, \dots, k_{r+1}\} \subseteq \mathbb{N}$ such that $x_1 = s + k_1a + 1, x_2 = s + k_1a + 1 + k_2a + 1, \dots, x_r = s + k_1a + 1 + \dots + k_ra + 1$ and thus $\text{next}_S(s) = s + k_1a + 1 + \dots + k_ra + 1 + k_{r+1}a + 1$. Therefore, $\text{next}_S(s) - s = k_1a + 1 + \dots + k_ra + 1 + k_{r+1}a + 1$. As $S \in E(a)$ then $\text{next}_S(s) - s - 1 = ta$ for some $t \in \mathbb{N}$. Consequently, $k_1a + 1 + \dots + k_ra + 1 + k_{r+1}a + 1 = ta + 1$. Then $(k_1 + \dots + k_{r+1})a + r + 1 = ta + 1$ and thus r is a multiple of a . Assume that $r = la$ for some $l \in \mathbb{N} \setminus \{0\}$

To conclude the proof, we check that $S \cup \{x_{(l-1)a+1}, x_{(l-1)a+2}, \dots, x_{(l-1)a+a}\} \in E(a)$. In order to see this, it is enough to see that $x_{(l-1)a+1} - s - 1$ is a multiple of a . This is true because $x_{(l-1)a+1} - s - 1 = s + k_1a + 1 + k_2a + 1 + \dots + k_{(l-1)a+1}a + 1 - s - 1 = (k_1 + \dots + k_{(l-1)a+1})a + (l-1)a + 1 - 1$ is a multiple of a . \square

Given a sequence of nonnegative integers $n_1 < n_2 < \dots < n_p$, we say that it is equidistant modulo a if $n_{i+1} - n_i - 1$ is a multiple of a for all $i \in \{1, \dots, p-1\}$.

Let S be a numerical semigroup. An element of $s \in S$ is called *a-refinable* if there exists $\{x_1 < x_2 < \dots < x_a\} \subseteq \{x \in \mathbb{N} \mid s < x < \text{next}_S(s) \text{ and } x_a < F(S)\}$ such that $S \cup \{x_1, x_2, \dots, x_a\}$ is a numerical semigroup and the sequence $s, x_1, x_2, \dots, x_a, \text{next}_S(s)$ is equidistant modulo a . We denote by $\mathcal{R}(S) = \{s \in S \mid S \text{ is } a\text{-refinable}\}$.

Theorem 4.6 *Let $S \in E(a)$ with $F(S) = F$. Then S is a maximal element in the set $\{T \in E(a) \mid F(T) = F\}$ if and only if $\mathcal{R}(S) = \emptyset$.*

Proof *Necessity.* If $\mathcal{R}(S) \neq \emptyset$, then there exists $s \in \mathcal{R}(S)$. Hence, there exist $\{x_1 < x_2 < \dots < x_a\} \subseteq \mathbb{N} \setminus \{0\}$ such that the sequence $s < x_1 < x_2 < \dots < x_a < \text{next}_S(s)$ is equidistant modulo a with $x_a < F$ and $S \cup \{x_1, x_2, \dots, x_a\}$ is a numerical semigroup. We deduce that $S \cup \{x_1, x_2, \dots, x_a\} \in E(a)$ with $F(S \cup \{x_1, x_2, \dots, x_a\}) = F$ contradicting the maximality of S .

Sufficiency. If we suppose that S is not maximal, then there exists $T \in \mathbf{E}(a)$ with $F(T) = F$ and $S \subsetneq T$. Let $x = \max(T \setminus S)$ and $s \in S$ such that $s < x < \text{next}_S(s)$. By applying Proposition 4.5, we obtain that $s \in \mathcal{R}(S)$ and thus $\mathcal{R}(S) \neq \emptyset$. \square

5. The elements of $\mathbf{E}(a)$ with maximal embedding dimension

Let S be a numerical semigroup. The cardinality of $\text{msg}(S)$ is known as the embedding dimension of S and it is denoted here by $e(S)$. In [8, Proposition 2.10] shows that $e(S) \leq m(S)$. We say that S has maximal embedding dimension (MED-semigroup) if $e(S) = m(S)$. The following result can be deduced from [2, Proposition I.2.9].

Lemma 5.1 *Let S be a numerical semigroup. Then S is a MED-semigroup if and only if $\{s - m(S) \mid s \in S \setminus \{0\}\}$ is a numerical semigroup.*

Proposition 5.2 *Let S be a MED-semigroup. Then $S \in \mathbf{E}(a)$ if and only if $T = \{s - m(S) \mid s \in S \setminus \{0\}\}$ is an element of $\mathbf{E}(a)$ and $m(S) \equiv 1 \pmod{a}$.*

Proof *Necessity.* By Lemmas 5.1 and 2.1, we have that T is a numerical semigroup and $m(S) \equiv 1 \pmod{a}$, respectively. To conclude the proof, it suffices to see that $\text{next}_T(t) - t - 1$ is a multiple of a for every $t \in T$. If $t \in T$, then there exists $s \in S \setminus \{0\}$ such that $t = s - m(S)$ and so $\text{next}_T(t) = \text{next}_S(t) - m(S)$. Therefore, $\text{next}_T(t) - t - 1 = \text{next}_S(s) - m(S) - (s - m(S)) - 1 = \text{next}_S(s) - s - 1$ is a multiple of a , because $S \in \mathbf{E}(a)$.

Sufficiency. Let us see that $S \in \mathbf{E}(a)$, that is, if $s \in S$ then $\text{next}_S(s) - s - 1$ is a multiple of a . If $s = 0$, then $\text{next}_S(s) - 0 - 1 = m(S) - 1$ is a multiple of a . If $s \neq 0$, then $s - m(S) = t \in T$ and $\text{next}_T(t) = \text{next}_S(s) - m(S)$. Hence, $\text{next}_S(s) - s - 1 = \text{next}_T(t) - t - 1$ is multiple of a , because $T \in \mathbf{E}(a)$. \square

From Lemma 5.1, it is easy to deduce the following result.

Lemma 5.3 *Let S be a numerical semigroup and $x \in S \setminus \{0\}$. Then $S(x) = (\{x\} + S) \cup \{0\}$ is a MED-semigroup with multiplicity x . Moreover, every MED-semigroup is of this form.*

Proposition 5.4 *Let $S \in \mathbf{E}(a)$ and $x \in S \setminus \{0\}$ such that $x \equiv 1 \pmod{a}$. Then $S(x) = (\{x\} + S) \cup \{0\}$ is an equidistant MED-semigroup modulo a . Moreover, every equidistant MED-semigroup modulo a is of this form.*

Proof By Lemma 5.3, we obtain that $S(x)$ is a MED-semigroup with multiplicity x . Clearly $S = \{s - x \mid s \in S(x) \setminus \{0\}\}$ and thus, applying Proposition 5.2, we obtain that $S(x)$ is equidistant modulo a .

Let T be an equidistant MED-semigroup modulo a . Then by Proposition 5.2, we get that $S = \{T - m(T) \mid t \in T \setminus \{0\}\} \in \mathbf{E}(a)$ and $m(T) \equiv 1 \pmod{a}$. Therefore, $T = (\{m(T)\} + S) \cup \{0\}$ with $S \in \mathbf{E}(a)$ and $m(T) \in S$ such that $m(T) \equiv 1 \pmod{a}$. \square

Let S be a numerical semigroup and let $n \in S \setminus \{0\}$. The Apéry set (named so in honour of [1]) of n in S is

$$\text{Ap}(S, n) := \{s \in S \mid s - n \notin S\}.$$

Lemma 5.5 [8, Lemma 2.4] *Let S be a numerical semigroup and $n \in S \setminus \{0\}$. Then $\text{Ap}(S, n) = \{0 = w(0), w(1), \dots, w(n - 1)\}$, where $w(i)$ is the least element in S congruent with i modulo n , for all $i \in \{0, \dots, n - 1\}$.*

Observe that the above lemma in particular implies that the cardinality of $\text{Ap}(S, n)$ is n . From [10], we can deduce the next result.

Proposition 5.6 *Let S be a numerical semigroup, $n \in S \setminus \{0\}$ and $T = (\{n\} + S) \cup \{0\}$. Then the following conditions hold:*

1. T is MED-semigroup
2. $m(T) = n$.
3. $F(T) = F(S) + n$.
4. $g(T) = g(S) + n - 1$.
5. $\text{msg}(T) = \text{Ap}(S, n) + \{n\}$.

Example 5.7 *It is clear that $S = \langle 5, 8, 9 \rangle$ is an equidistant numerical semigroup modulo 2. We have that 9 is an element in S such that $9 \equiv 1 \pmod{2}$. Then from Proposition 5.4, we obtain that $T = (\{9\} + S) \cup \{0\}$ is an equidistant MED-semigroup modulo 2. Since $F(S) = 12$, $g(S) = 8$ and $\text{Ap}(S, 9) = \{0, 5, 8, 10, 13, 15, 16, 20, 21\}$, by Proposition 5.6 we have that $F(T) = 12 + 9 = 21$, $g(T) = 8 + 9 - 1 = 16$ and $\text{msg}(T) = \{9, 14, 17, 19, 22, 24, 25, 29, 30\}$.*

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